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Chapter 14

Split of Composite Components for Distributed Applications

Ansgar Radermacher, Önder Gürcan, Arnaud Cuccuru, Sébastien Gérard, and Brahim Hamid

Abstract Composite structures as in UML are a way to ease the development of complex applications. Composite classes contain sub-components that are instantiated, interconnected and configured along with the composite. Composites may also contain operations and further attributes. Their deployment on distributed platforms is not trivial, since their sub-components might be allocated to different computing nodes. In this case, the deployment implies a split of the composite. In this chapter, we will motivate why composites need to be allocated to different nodes in some cases by examining the particular case of interaction components. We will also discuss several options to achieve the separation and their advantages and disadvantages including modeling restrictions for the classes.

14.1 Introduction

The basic idea behind any component-oriented approach is that elementary application pieces (i.e. components) can be composed together in order to achieve the functionality of a more complex system. Component-oriented approaches are usually grounded on a design process including component development or reuse, assembly and deployment.

In the **component assembly** step, the system under design is itself considered as a component. It is hierarchically defined by an assembly of existing components using an Architecture Description Language (ADL) [4], where the assembly is concretely specified by connections expressed between sub-components (parts). In the context of this chapter, we focus on Unified Modeling Language (UML) [16] as modeling language. Sub-components can themselves be defined as assemblies, resulting in hierarchical systems of arbitrary depth.

In the **deployment specification** step, the target execution platform for the application is considered. The model of the execution platform usually consists, at least, of an identification of the various execution nodes, as well as available communication paths between them. The deployment specification consists of allocating the components of the application model to execution nodes of the platform (often indirectly by allocating them to processes or threads which in turn are allocated

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to execution nodes, but we simplify this aspect in the context of this chapter). Allocation is usually done taking into account *non-functional* requirements of the system under design, such as execution time constraints, memory footprint, communication throughput, etc.

It is sometimes necessary to allocate sub-components to different execution nodes which requires a split of the associated composite. The next section illustrates this problem by means of a small example, section 14.3 provides multiple options how to split composites. Section 14.4 examines how existing component frameworks split composites. An evaluation and comparison of these options is given in section 14.5. Section 14.6 concludes this chapter.

14.2 Motivating Example

In this section, we motivate why some composites need to be split by examining interaction components.

Consider a very simple application with two components, **A** and **B** as shown in Fig. 14.1. **A** has a port **q** with a required interface **I**, **B** has a port **p** with a provided interface **I**.

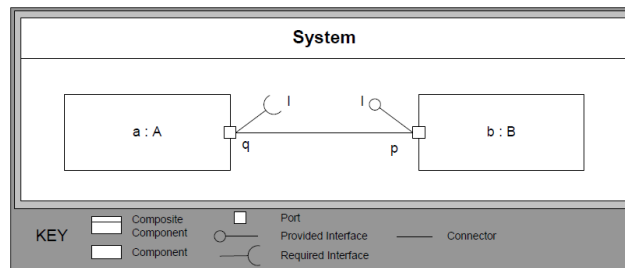


Fig. 14.1 A simple system with two components and uni-directional communication.

Now consider that the communication between **A** and **B** is realized by a component that implements the interaction on top of the operating system's socket API. We call such a component an interaction component (also called connector in the context of the DDS-for-CCM specification [15]). On a logical level, this component is a single entity that may contain configuration data such as a port number, connection policies or a unique identifier (object reference).

If we want to distribute the application onto two nodes, **a** and **b** are allocated to different nodes. Fig. 14.2 shows the architecture of the example system. Please note that the composite structure diagram distinguishes between a role (corresponding to a kind of instance) and its type, i.e. the socket is not a nested classifier within the system but a part of the system on an instance level. Thus, the first component that is split is the component representing the system itself (**System**). However, the **System** component is a particular case, since there exists only one instance, it has no behavior of its own and there are no connections from the system boundary to inner parts (called delegation connectors in UML [16]). Thus, it is a pure assembly component and basically used to define the instances of a system and their interconnections.

Fig. 14.3 shows the internal structure of the **SocketConnector** component. It consists of a client and server stub (**cli** and **srv** respectively) which both access a socket run-time. The dashed outline of the latter indicates that this component is shared: it is not instantiated along with **SocketConnector** but exists independently. The access to a shared resource within a composite corresponds to a

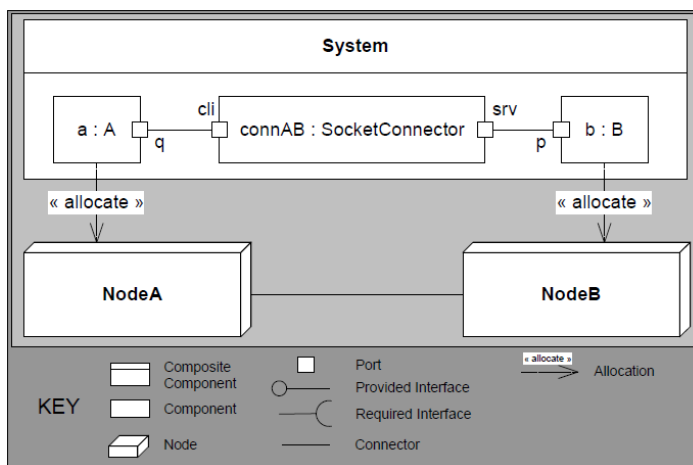


Fig. 14.2 Distribution of the system given in Fig. 14.1. The System composite component is now with sockets and allocation.

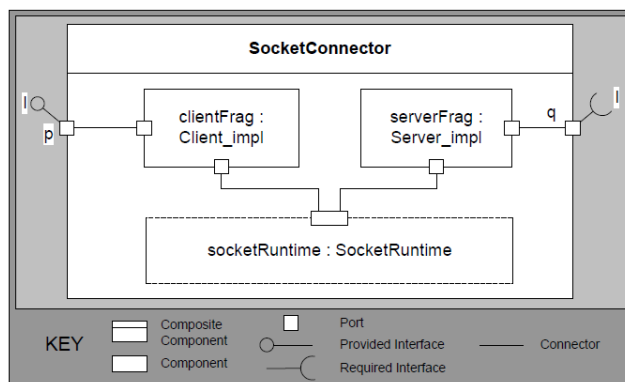


Fig. 14.3 Internal structure of the SocketConnector composite component given in Fig. 14.2.

kind of vertical connection: the communication of the stubs with the run-time is a communication between different layers, pre-assembled within the composite.

Since the communication with the interaction component is a simple local communication, the interaction component itself needs to be separated. We can further follow local connections within the connector to determine the allocation of the internal parts of the connector. The allocations within the socket connector can thus be derived from the allocations of the application components: the client fragment of the connector needs to be co-located with A and the server fragment with B. An interesting aspect is the socket run-time that is shared by client and by server fragment. Whereas it exists only once from a logical viewpoint, it must be present on each node and thus be allocated to NodeA and to NodeB. Fig. 14.4 shows the resulting split of the SocketConnector.

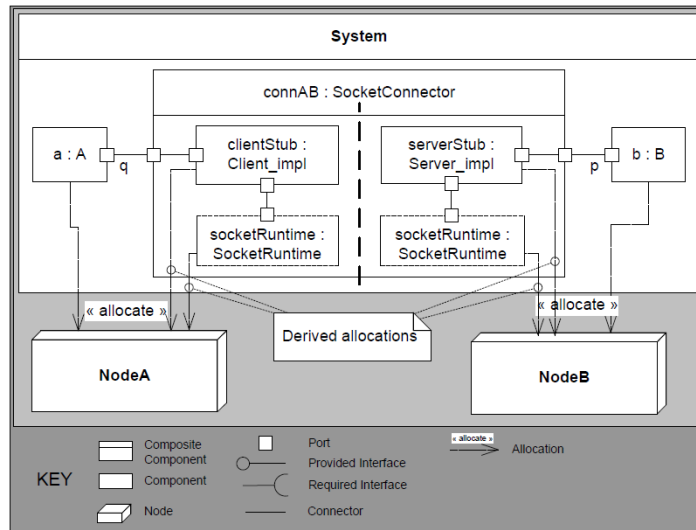


Fig. 14.4 Splitting the composite component `SocketConnector` during the distribution of the system.

Since a composite can enable distribution, its split should be authorized under the condition that *this split does not modify the component's semantics*¹. This is the case, if a composite does not have a behavior of its own (only delegation to parts), nor any configuration data. Since the latter is too strict, the composite may offer virtual configuration attributes that are effectively realized by its parts. This means that the configuration attributes of the composite are linked with configuration attributes in the parts. The same attribute might appear in multiple parts.

Now consider a slight extension of the example: `B` also talks to `A`, using the same interface, `A` has an additional port `p`, `B` an additional port `q` and both are connected, as shown in Fig. 14.5.

In this case two parts (`connAB` and `connBA`) are typed with `SocketConnector`. But, the allocation of the sub-part is different for the two instances (parts):

Since `a` is on `NodeA`, the `clientStub` part of the instance `connAB` must be on `NodeA` as well to satisfy the co-localization constraint caused by the assumption of inseparable simple connections. But with the same argument, `clientStub` of instance `connBA` must be on `NodeB`, co-localized with `b`. Thus, allocation is instance based and it might happen that two different instances of a composite have different allocation specifications for their parts. Thus, the split is not trivial and we will study multiple options how to split the composite in case of the example in section 14.3.

¹ Preserving semantics of components is also important in order to be able to analyze them correctly [11].

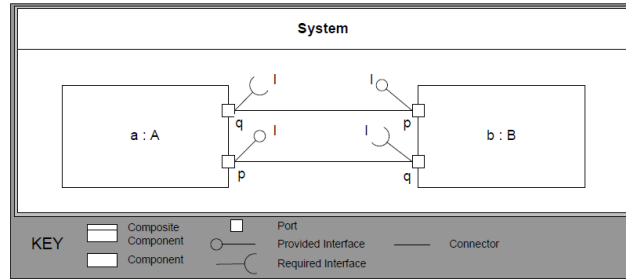


Fig. 14.5 A simple system with two components and bi-directional communication.

14.3 Different ways to split composites

In the sequel, three different options to split composites are shown by means of the simple example from section 14.2.

Option 1 – Keep composites

The first option is to keep a modified variant of the composite that only contains the subset of parts which are deployed on a certain node. Fig. 14.6 shows the result for the uni-directional variant of the example: 'SocketConnector' is the variant of the original 'SocketConnector'. It contains the subset of parts that are allocated on NodeA, 'clientStub' and 'socketRuntime'. Note that splitting is in general not trivial, since the split must also consider super-classes. In our case, the ports of the socket are inherited by an abstract interaction component (aka connector type). Depending on how the super-class is organized, the composite only inherits from a subset of super-classes or super-classes need to be split as well which complicates the design.

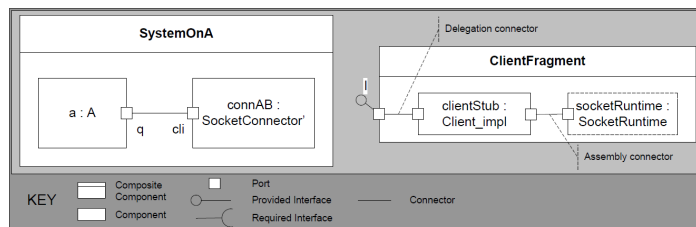


Fig. 14.6 Option 1: Splitting components in the uni-directional example given in Fig. 14.1 by keeping composites.

Please note that it is not the part that is allocated on a certain node, but the (sub-) instance that is associated with a part. If there is a second instance whose sub-instances are allocated in a different way, a second variant of the composite with a part subset must be created. This is shown in Fig.

14.7. The creation of multiple variants implies a certain overhead which –although small– may be non-acceptable on resource constraints systems.

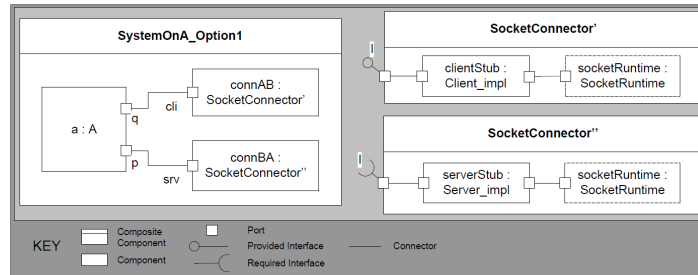


Fig. 14.7 Option 1: Splitting components in the bi-directional example given in Fig. 14.5 by *keeping* composites.

Option 2 – Flatten composite

Flattening a composite component is a well-known approach in the literature [6, 7, 11] in which a composite component may disappear in the deployment model, i.e. it is replaced by its internal structure. The internal assembly connections of a composite become assembly connections of the containing composite (the System class in case of the example). The delegation connections² refine the final targets of existing assembly connections in the containing composite.

Fig. 14.8 shows the example system for *NodeA*, in which the *SocketConnector* composite component has been flattened. The two parts in the system typed by a socket implementation have been replaced by parts that are directly typed with elements of the socket implementation. The original composition hierarchy may still be visible via a suitable naming convention for these new parts by prefixing them with the original part name, as done in the example with the prefixes *connAB* and *connBA*.

Option 3 – Flatten composite, require explicit fragment sub-components

The third option is a variation of the second solution. We also flatten the *SocketConnector* composite component, but require that the composite must contain exclusively specific sub-components that we call *fragments*. A fragment encapsulates the parts of a composite that are allocated on the same node, conversely each fragment within a composite is typically allocated on a different node. The latter implies a restriction that is verified by a validation rule: fragments may not be connected by UML assembly connectors. The modeling of *SocketConnector* with fragments is shown in Fig. 14.9.

The resulting system is shown in Fig. 14.10. The composite has been flattened; the fragments have become top-level elements. The result looks very similar as the solution in Fig. 14.7, effectively the explicitly modeled fragments replace the derived subsets of the composite.

² *Assembly connectors* are connections between inner parts, *delegation connectors* connections from the composite to an inner part.

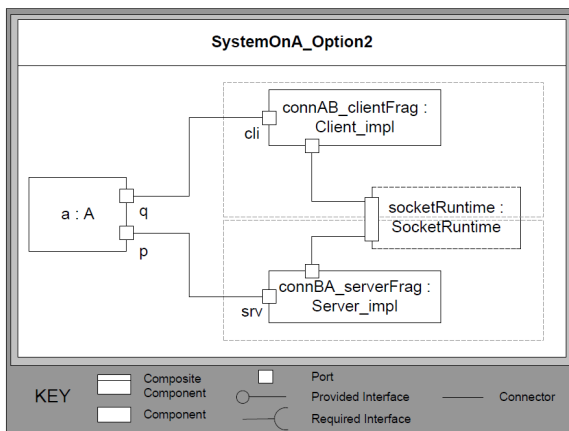


Fig. 14.8 Option 2: Splitting components in the bi-directional example given in Fig. 14.5 by flattening composites.

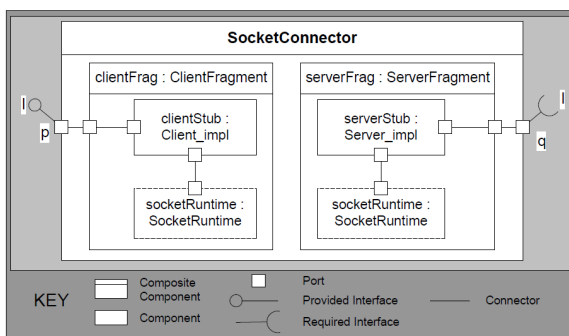


Fig. 14.9 Option 3: The SocketConnector composite component with explicit fragments.

14.4 Support for splitting composites in existing frameworks

In the following, we sketch the existing component frameworks that have a specific support for interaction components³ and show how these frameworks may handle composite splitting, mainly in the context of interaction components.

³ Having specific support for interaction components is needed in order to be able to address the composite split in a systematic way.

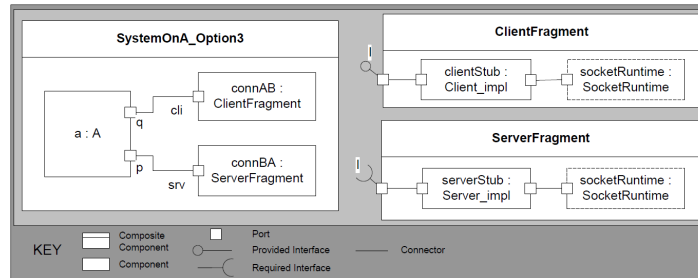


Fig. 14.10 Option 3: Splitting components in the bi-directional example given in Fig. 14.5 by flattening the System composite with explicit fragments.

14.4.1 DDS for CCM

The connector element that we have used in the motivating example is supported in multiple component models. As already mentioned, it has been standardized within the context of the OMG (Object Management Group) standard CCM (CORBA component model) [13]. More specifically, it is part of the DDS for CCM [15] specification, enabling component interactions via OMG's Data Distribution Service. Within this specification, the term GIS (Generic Interaction Support) is introduced. GIS will be part of the upcoming OMG unified component model [17]. The underlying connector extension for CCM has been proposed in [19]. Deployment with CCM is based on the specification for Deployment and Configuration (D&C) of distributed component-based applications [14]. The D&C standard describes a so-called deployment plan, a specification of instances that refer to component implementations, the interconnections between these instances, their configuration and their allocation to a node.

In the DDS for CCM specification, DDS interaction components are not identified as composites, since there are separate writer and consumer components. This is useful in case of DDS in which connections are implicitly created by sharing the same topic, i.e. there is no single component that represents an interaction. But the generic interaction support enables explicit point-to-point interactions for which composites would be useful. D&C supports two kinds of implementations of software components [14]:

- Monolithic implementations, where the code of the composite component is compiled as a single block.
- Assembly (composite) implementations, including the set of implementation of all the parts that the composite component includes. There must eventually be monolithic implementations at the “leaves” of the hierarchical implementation. Assembly allows dependent packages to be deployed on distinct target nodes, enabling flexibility in composite component instantiation.

While the D&C specification allows composites, the composites have no identity and cannot be reused. This has been analyzed in [10]. In this article, the authors review and compare the ability of 13 component models of handling component composition. They identify the development with D&C as a “deposit only” repository for composites: a composite component that results from the component assembly step can be deposited in a repository but cannot be retrieved from it, because it does not have an identity of its own. In the end, only monolithic components are deployed, i.e. the component hierarchy is flattened. Note that this does not only apply to interaction components but to all composite components, even if they are deployed on the same node, i.e. a stronger variant of the *flatten* option in section 14.3.

14.4.2 *Fractal*

Connectors have also been introduced in the context of Fractal⁴ [2, 5]: a binding is defined as a communication path between component interfaces. Bindings can be *primitive* or *composite*. A primitive binding (direct connector) binds one client interface and one server interface in the same address space. A composite binding is a communication path between an arbitrary number of distributed component interfaces and is represented as a set of primitive bindings and binding components. *Binding components* are called Fractal connectors, and are normal Fractal components, whose role is dedicated to communication.

However, there is no support for splitting in Fractal in the sense of interaction components as shown in Section 14.2.

14.4.3 *SOFA*

In SOFA⁵ [1, 3, 12], connectors are used to support transparent distribution of applications. A connector might support a transport mechanism such as CORBA⁶ or low-level socket mechanisms. In this context, they are responsible for marshalling and unmarshalling and interfacing with the transport layer. But they can also be used for synchronization or interception. Connectors are automatically generated.

In SOFA, the connector plugging is performed after component instantiation using a split of the connector into two parts: the *server* and the *client* connector units (fragments). Whenever component interfaces query a connector reference, the corresponding *server* connector unit is returned (instead of returning a reference directly to an interface). Similarly, whenever an interface is being connected to another component, a *client* connector unit is created and bound. The connector composite specifies the parts, into which it is later split explicitly, corresponding to the *fragment* option (Option 3).

14.4.4 *Qompass*

The FCM [8] (Flex-eWare component model) component model has the objective to unify the component models of Fractal and CCM. It extends the UML composite structures with dedicated interaction components – as, for instance, the socket connector presented in the motivating example (Section 14.2) – flexible ports and container services. This component model is supported by an add-on to the Papyrus⁷ UML modeler called Qompass designer. Fig. 14.11 shows the definition of the socket interaction component within a Qompass modeling library and the Qompass context menu. This add-on was first introduced as eC3M (embedded Component Container Connector Model/Middleware) [18]. Upon deployment, the tool chain executes a model transformation that replaces annotated UML connectors with the associated interaction components, as shown in the motivating example given in Section 14.2 (the transition from Fig. 14.1 to Fig. 14.2). This transformation includes an instantiation of the interaction component to the context in which it is used (similar to the generation of in SOFA). A further model transformation produces a model per

⁴ The Fractal Component Model, <http://fractal.objectweb.org/specification/>, last access on 07/02/2014.

⁵ SOFA 2, <http://sofa.ow2.org/>, last access on 07/02/2014.

⁶ Common Object Request Broker Architecture, <http://www.corba.org/>, last access on 07/02/2014.

⁷ The Papyrus UML modeler – <http://www.eclipse.org/papyrus>, last access on 23.01.2014

node. During the latter, the composites within the FCM models are split. The composites that are concerned are mainly interaction components and the dedicated system component.

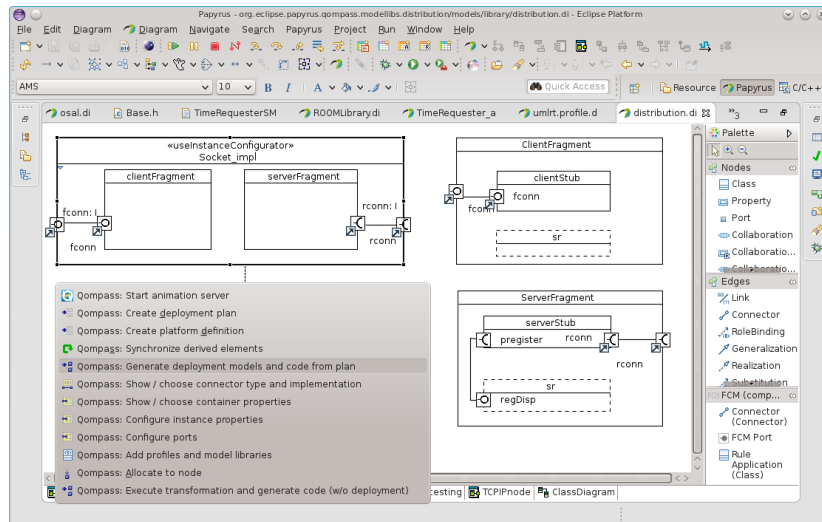


Fig. 14.11 Definition of a socket connector within Papyrus, using the client and server *fragments*

In Qompass, interaction components with explicitly identified fragments are flattened, i.e. the *fragment* option (Option 3). Being based on UML, Qompass must handle the specific case of a dedicated system component. Such a component is required, since connections can only be defined in the context of an enclosing composite (unlike for instance in D&C). Thus, Qompass must also split the system component, if the contained components are deployed on different nodes. The approach that has been chosen is to create a specific variant of the *System* component on each node, i.e. the *keep* option (Option 1). Note that it is not possible to *flatten* the system component, since the UML component model requires an enclosing composite for defining connections.

14.5 Discussion

Obviously, all splitting options increase the number of classes. When the composites are *kept* (Option 1), there is no need to remove additional assembly connections from the system. *Flattening* (Option 2) makes top-level composites bigger, since these composites have to incorporate the contents of a flattened component (sub-components and their connections) instead of the component itself. In *fragmentation* (Option 3), a possible split is anticipated and explicitly defined by the developer. Since the composite may not have assembly connectors, no additional connectors are added to the *System* class (the composite that contains the split composite). Based on this observation, to make a quantitative comparison, we measured the *footprints* associated with the different splitting options. The code size of a complete application has been measured in case of the simple uni-directional system and its bi-directional variant for splitting options 1, 2 and 3, as shown in Table 14.5. The

results were obtained on a Linux machine with gcc 4.7 (optimizations disabled). As expected, *flattening* (Option 2) results in a slightly smaller footprint compared to the other two.

Table 14.1 Footprint of different deployment options for the simple uni-directional (Fig. 14.1) and simple bi-directional (Fig. 14.5) systems. The results show that flattening has the smallest footprint among three deployment (splitting) options.

example system	code size (bytes)		
	Opt 1 (keep)	Opt 2 (flatten)	Opt 3 (fragments)
Simple uni-direct	13904	12233	13936
Simple bi-direct	14668	13754	14710

However, *flattening* is evidently not possible for a top-level component, since the transformation towards a model having only monolithic components and assembly connections⁸ is rather straight forward. Thus, the resulting system is different, since the internal connections become visible in the system. This may be annoying, if the same composite is instantiated more than once in the original model, e.g., if we have more than one socket connector. Also note that the internal structure of an interaction component might be more complex than the simplified `SocketConnector` used for illustration purposes. This makes it a bit difficult to link it with the original model, for instance when debugging is done on the level of the deployed model, but fixes must be made in the original design model. Other tasks that are affected by this difference are for instance trace mechanisms (which must translate a trace specification for a composite into suitable specifications for the inner parts) and the replacement of a composite implementation with another one (e.g. in the context of different system configurations). The advantage is a slightly reduced footprint and a resolution for the splitting problem.

Another important thing to consider is debugging. *Debugging* is generally defined as the process of locating and fixing or bypassing bugs in the underlying software, to achieve reliable systems. To this end, various debugging tools are developed help to identify errors at the various stages of the software development process. Especially, debugging and visualizing the behavior of component-based embedded software using models such as the Unified Modeling Language (UML) [16] diagrams has become a reality. For instance, model-based tools such as Papyrus and the commercial tool Rhapsody⁹ (“live-animation” features) enable model-based debugging of embedded software systems using sequence diagrams and state charts. In case of Papyrus, animation is based on an injected probe that communicates with the development environment. Showing the activation of a delegation connector within a composite is evidently only possible, if the composition hierarchy has not been flattened. Hence, the closer the deployed model is to the original architecture, the easier it is to debug. In this sense, since *keeping* the original composition hierarchy (Option 1) has the advantage that the deployed model is closer to the original architecture, it is a bit easier to debug compared to Option 2 and 3.

Besides, in some domains (such as aerospace and electrical cars) the overall architecture of vehicles becomes very complex [20]. One possibility to tackle this complexity at run-time, is the use of dynamic reconfiguration abilities. *Dynamic reconfiguration* is a process of modifying the software architecture and enact the modifications during the system’s execution [9], which means making the software evolve from one configuration to another at run-time, as opposed to design-time, while introducing little or ideally no impact on the systems execution. This prevents the system to be taken

⁸ In UML-like languages, connectors are always owned by a composite, i.e. a System composite must be kept.

⁹ IBM Rational Rhapsody Developer, <http://www-03.ibm.com/software/products/en/ratirhap/>, last access on 04.02.2014

off-line and/or restarted to accommodate changes. Considering the split of composite components discussed in this chapter, a dynamic reconfiguration would replace the `SocketConnector` component with another interaction component. In order to be able to do this seamlessly, splitting Option 1 is better suited since we do not need to remove additional assembly connections from the system.

14.6 Conclusion

We have shown that the deployment of composite instances which are partly allocated on one node and partly on another can be tackled in several ways with different advantages and disadvantages. The choice of a suitable split option depends on properties of the composite that should be split. For instance, in Qompass designer we keep the composite of the `System` component, since this particular component (no inheritance, single instance) can be split easily and since *flattening* would result in multiple top-level components. On the other hand, we flatten interaction components and require the explicit use of fragments, since we want to avoid the problems that come with multiple instances (creating potentially multiple variants of a split component). The choice depends also on the deployment goals, e.g. whether an optimized application compared to a debug-enabled application should be delivered. The options are rather evident, but –to our knowledge– the task had not been examined systematically earlier.

The interest of deploying composites with complex allocation properties is not artificial: a composite definition is a suitable choice for interaction components enabling distribution. In this context, the raised issues concern principally framework and tool developers, i.e. developers of interaction components and developers of model transformations associated with the split of composites. However, the results also apply to a sub-system modeled by composite classes that need to be allocated on multiple execution nodes. In this case, system modellers or designers are concerned since they need to respect restrictions associated with the split of a composite and should know the consequences of different split options.

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